

True or False: (we only address the **early** FORTRAN covered in class)

1. FORTRAN has the **recursive calls** feature where SUBROUTINES & FUNCTIONS can call themselves recursively, for execution time power and efficiency. **F**
(recursion leads to dynamic memory de/allocation of AR's in the stack, involving OS)
2. The compiler will generate the code that will build the subroutines/functions calls' activation records (ARs) later at run time. **T**
3. The FORTRAN COMMON blocks facility is mainly an ad-hoc controlled **virtual** global block-name visibility to save/share memory space at run time. **T**
4. In FORTRAN, any declared name outside the main program is **global** in visibility. **F**
Only SUB's/FUNC's names are global!
5. FORTRAN would be **secure**, the moment it adopts the explicit name declaration. **F**
NOP! Still there are other features that makes the language insecure, e.g., overworking types!
6. In a modern block-structured HLL where the contour diagram definition is used to define name visibility, assume subroutine **A** declares subroutine **S** which declares function **F**, the names **S&F** are considered part of **A**'s locally declared names and will be seen everywhere within **A**'s contour box. **F**
Only **S** will be visible inside **A**'s box, never **F**, since it is declared inside **S**, and **S**'s box is one way mirror that never lets a name declaration (**F**) seeps to the outside!
7. According to the contour diagram rules, an id-name **X** might have different meanings (declarations) in the program code (nested box's). Its meaning will depend on where we use it! **T**
Yes, statically scoped block-structured HLL's allows the same name to have different declaration in different box's in the program!
8. According to the contour diagram definition, the compiler will report an "undeclared name/id X" error when the compiler encounters the use of **X** outside the **scope** of its declaration. **T**
9. FORTRAN is a clear example of a language that strived for maximum **efficiency & power** (memory and execution time) while sacrificing **security**. **F**
Gone all the way for **efficiency** but not **power**!
10. The activation record (AR) of the main program in any block-structured HLL must be the first AR to be generated at run time in the memory. **T**

11. The *caller's* return address will be save in its own AR. T
12. The dynamic link (DL) is a pointer from the *callee's* AR to the *caller's* AR. T
13. The use of subroutines/functions in a program is a tradeoff between ***abstraction*** (with all associated advantages) and the dynamic ***overhead*** on the increase of the host code execution time. T
14. FORTRAN is *dynamically* type checked. F (No way! We strive for efficiency!)
15. FORTRAN uses ***stack*** to de/allocate subroutines/functions' AR's at run time. F (No way! We strive for efficiency! Stack will force dynamic de/allocation of memeory)