

True or False: (questions are about FORTRAN & Algol, early versions , ONLY)

- a) Algol is more powerful than FORTRAN. T  
Algol has Dynamic arrays, by-name, rich typing system, etc.
- b) Upon calling a procedure/function that does not define any local names, an activation record (AR) will not be pushed into the runtime stack, in such special case. F  
NOP! An AR is created regardless, for other duties as explained in class!
- c) “*Recursion*” and *dynamic arrays* in Algol are examples of gaining **power** for loss of **security**. F  
NOP! power versus efficiency, not security, both mechanisms are secure in Algol.
- d) Algol is a *block-structured* imperative High Level Language. T
- e) Unlike FORTRAN, Algol designers NEVER trusted the users to do the right thing, hence the language is secure! F  
NOP! The trusted users to do the correct by-name parameter passing!
- f) Algol's *dynamic arrays* alleviated the problem of having a lack of *pointers* to facilitate the implementation of dynamic data structures, e.g., lists, stacks, and queues. T  
The use of arrays, dynamic arrays, pointers in dynamic ADTs is explained in class.
- g) Algol's typing system is *insecure* due to the use of *by-name* parameter passing. F  
By-name is not related to the typing system!
- h) The *by-name* parameter passing facility in Algol is an examples of gaining **power** and losing **security** and **efficiency**. T
- i) In Algol, *nested blocks* inside *procedures* might be called from any place in the code, yet according to the contour diagram rules in statically scoped languages. F  
Blocks are NEVER called, they are only in-place static abstraction.
- j) In Algol, the *type* declaration section is *elaborated* at run time, for better *efficiency*. F

**Extra credits** (if correct you **get +5 points** for each correct answer, otherwise you **lose -5 points**):

- k) The *insecurity* of Algol's *by-name* will be solved if *thunks* are evaluated **only once**, at the caller environment, which won't affect the *by-name*'s **power**. F  
NOP! It will violate its semantics & integrity since it won't be powerful anymore!
- l) In Algol, in some special scenarios, the AR of the *caller* as well as the *definer* of any invoked (called) procedure **P** (callee) are **NOT** always guaranteed to be pushed/exist in the run time stack when **P** is called, i.e., below **P**'s AR. F  
NO way, not in statically scoped and stack model of computation HLLs, like Algol!